

Programming Assignment III

Due Tuesday, April 7, 2020 (19 Farvardin)

1 Introduction

In this assignment you will write a parser for Cool. The assignment makes use of two tools: the parser generator (the C++ tool is called `bison`; the Java tool is called `CUP`) and a package for manipulating trees. The output of your parser will be an abstract syntax tree (AST). You will construct this AST using semantic actions of the parser generator.

You certainly will need to refer to the syntactic structure of Cool, found in Figure 1 of The Cool Reference Manual (as well as other parts). The documentation for `bison` and `CUP` is available online. The C++ version of the tree package is described in the *Tour of Cool Support Code* handout. The documentation for the Java version is available online. You will need the tree package information for this and future assignments.

There is a lot of information in this handout, and you need to know most of it to write a working parser. *Please read the handout thoroughly.*

You must work individually for this assignment.

2 Files and Directories

To get started, create a directory where you want to do the assignment and execute one of the following commands *in that directory*. For the C++ version of the assignment, you should type

```
gmake -f /usr/class/cs143/assignments/PA3/Makefile
```

For Java, type:

```
gmake -f /usr/class/cs143/assignments/PA3J/Makefile
```

(notice the “J” in the path name). This command will copy a number of files to your directory. Some of the files will be copied read-only (using symbolic links). You should not edit these files. In fact, if you make and modify private copies of these files, you may find it impossible to complete the assignment. See the instructions in the README file. The files that you will need to modify are:

- `cool.y` (in the C++ version) / `cool.cup` (in the Java version)

This file contains a start towards a parser description for Cool. The declaration section is mostly complete, but you will need to add additional type declarations for new nonterminals you introduce. We have given you names and type declarations for the terminals. You might also need to add precedence declarations. The rule section, however, is rather incomplete. We have provided some parts of some rules. You should not need to modify this code to get a working solution, but you are welcome to if you like. However, do not assume that any particular rule is complete.

- `good.cl` and `bad.cl`

These files test a few features of the grammar. You should add tests to ensure that `good.cl` exercises every legal construction of the grammar and that `bad.cl` exercises as many types of parsing errors as possible in a single file. Explain your tests in these files and put any overall comments in the README file.

- README

As usual, this file will contain the write-up for your assignment. Your report must contain the following parts: First, explain your goals, design decisions and overview of your algorithm. Second, explain parts of your project such as functions, global variables,... . Third, bring an example and explain your code's output for it. It is part of the assignment to explain things in text, as well as to comment your code.

To submit your assignment, run `"/usr/class/cs143/bin/submit"` from your PA3 directory and follow the instructions.

3 Testing the Parser

You will need a working scanner to test the parser. You may use either your own scanner or the `coolc` scanner. By default, the `coolc` scanner is used; to change this behavior, replace the `lexer` executable (which is a symbolic link in your project directory) with your own scanner. Don't automatically assume that the scanner (whichever one you use!) is bug free—latent bugs in the scanner may cause mysterious problems in the parser.

You will run your parser using `myparser`, a shell script that “glues” together the parser with the scanner. Note that `myparser` takes a `-p` flag for debugging the parser; using this flag causes lots of information about what the parser is doing to be printed on `stdout`. Both `bison` and `CUP` produce a human-readable dump of the LALR(1) parsing tables in the `cool.output` file. Examining this dump is frequently useful for debugging the parser definition.

You should test this compiler on both good and bad inputs to see if everything is working. Remember, bugs in your parser may manifest themselves anywhere.

Your parser will be graded using our lexical analyzer. Thus, even if you do most of the work using your own scanner you should test your parser with the `coolc` scanner before turning in the assignment.

4 Parser Output

Your semantic actions should build an AST. The root (and only the root) of the AST should be of type `program`. For programs that parse successfully, the output of `parser` is a listing of the AST.

For programs that have errors, the output is the error messages of the parser. We have supplied you with an error reporting routine that prints error messages in a standard format; please do not modify it. You should not invoke this routine directly in the semantic actions; `bison/CUP` automatically invokes it when a problem is detected.

Your parser need only work for programs contained in a single file—don't worry about compiling multiple files.

5 Error Handling

You should use the `error` pseudo-nonterminal to add error handling capabilities in the parser. The purpose of `error` is to permit the parser to continue after some anticipated error. It is not a panacea and the parser may become completely confused. See the `bison/CUP` documentation for how best to use `error`. **In your README, describe which errors you attempt to catch.** Your test file `bad.cl` should have some instances that illustrate the errors from which your parser can recover. To receive full credit, your parser should recover in at least the following situations:

- If there is an error in a class definition but the class is terminated properly and the next class is syntactically correct, the parser should be able to restart at the next class definition.
- Similarly, the parser should recover from errors in features (going on to the next feature), a `let` binding (going on to the next variable), and an expression inside a `{...}` block.

Do not be overly concerned about the the line numbers that appear in the error messages your parser generates. If your parser is working correctly, the line number will generally be the line where the error occurred. For erroneous constructs broken across multiple lines, the line number will probably be the last line of the construct.

6 The Tree Package

There is an extensive discussion of the C++ version of the tree package for Cool abstract syntax trees in the *Tour* section of the Cool documentation. The documentation for the Java version is available on the course web page. You will need most of that information to write a working parser.

7 Remarks

You may use precedence declarations, but only for expressions. Do not use precedence declarations blindly (i.e., do not respond to a shift-reduce conflict in your grammar by adding precedence rules until it goes away).

The Cool `let` construct introduces an ambiguity into the language (try to construct an example if you are not convinced). The manual resolves the ambiguity by saying that a `let` expression extends as far to the right as possible. The ambiguity will show up in your parser as a shift-reduce conflict involving the productions for `let`.

This problem has a simple, but slightly obscure, solution. We will not tell you exactly how to solve it, but we will give you a strong hint. In `coolc`, we implemented the resolution of the `let` shift-reduce conflict by using a `bison/CUP` feature that allows precedence to be associated with productions (not just operators). See the `bison/CUP` documentation for information on how to use this feature.

Since the `mycoolc` compiler uses pipes to communicate from one stage to the next, any extraneous characters produced by the parser can cause errors; in particular, the semantic analyzer may not be able to parse the AST your parser produces.

8 Notes for the C++ version of the assignment

If you are working on the Java version, skip to the following section.

- Running `bison` on the initial skeleton file will produce some warnings about “useless nonterminals” and “useless rules”. This is because some of the nonterminals and rules will never be used, but these *should* go away when your parser is complete.
- You must declare `bison` “types” for your non-terminals and terminals that have attributes. For example, in the skeleton `cool.y` is the declaration:

```
%type <program> program
```

This declaration says that the non-terminal `program` has type `<program>`. The use of the word “type” is misleading here; what it really means is that the attribute for the non-terminal `program` is stored in the `program` member of the union declaration in `cool.y`, which has type `Program`. By specifying the type

```
%type <member_name> X Y Z ...
```

you instruct `bison` that the attributes of non-terminals (or terminals) `X`, `Y`, and `Z` have a type appropriate for the member `member_name` of the union.

All the union members and their types have similar names by design. It is a coincidence in the example above that the non-terminal

It is critical that you declare the correct types for the attributes of grammar symbols; failure to do so virtually guarantees that your parser won’t work. You do not need to declare types for symbols of your grammar that do not have attributes.

The `g++` type checker complains if you use the tree constructors with the wrong type parameters. If you ignore the warnings, your program may crash when the constructor notices that it is being used incorrectly. Moreover, `bison` may complain if you make type errors. Heed any warnings. Don’t be surprised if your program crashes when `bison` or `g++` give warning messages.

9 Notes for the Java version of the assignment

If you are working on the C++ version, skip this section.

- You must declare `CUP` “types” for your non-terminals and terminals that have attributes. For example, in the skeleton `cool.cup` is the declaration:

```
nonterminal program program;
```

This declaration says that the non-terminal `program` has type `program`.

It is critical that you declare the correct types for the attributes of grammar symbols; failure to do so virtually guarantees that your parser won’t work. You do not need to declare types for symbols of your grammar that do not have attributes.

The `javac` type checker complains if you use the tree constructors with the wrong type parameters. If you fix the errors with frivolous casts, your program may throw an exception when the constructor notices that it is being used incorrectly. Moreover, `CUP` may complain if you make type errors.

10 Extra credit

Cool is based on a single inheritance model. As a bonus activity, you may choose to continue the implementation of Cool on a multiple inheritance basis. The requirement for this alternation of basis to the syntactic structure of Cool is the simple substitution of the definition of class with the following:

```
class ::= class TYPE [inherits TYPE  $[[, TYPE]^*$ ] {  $[[feature;]^*$ }
```